

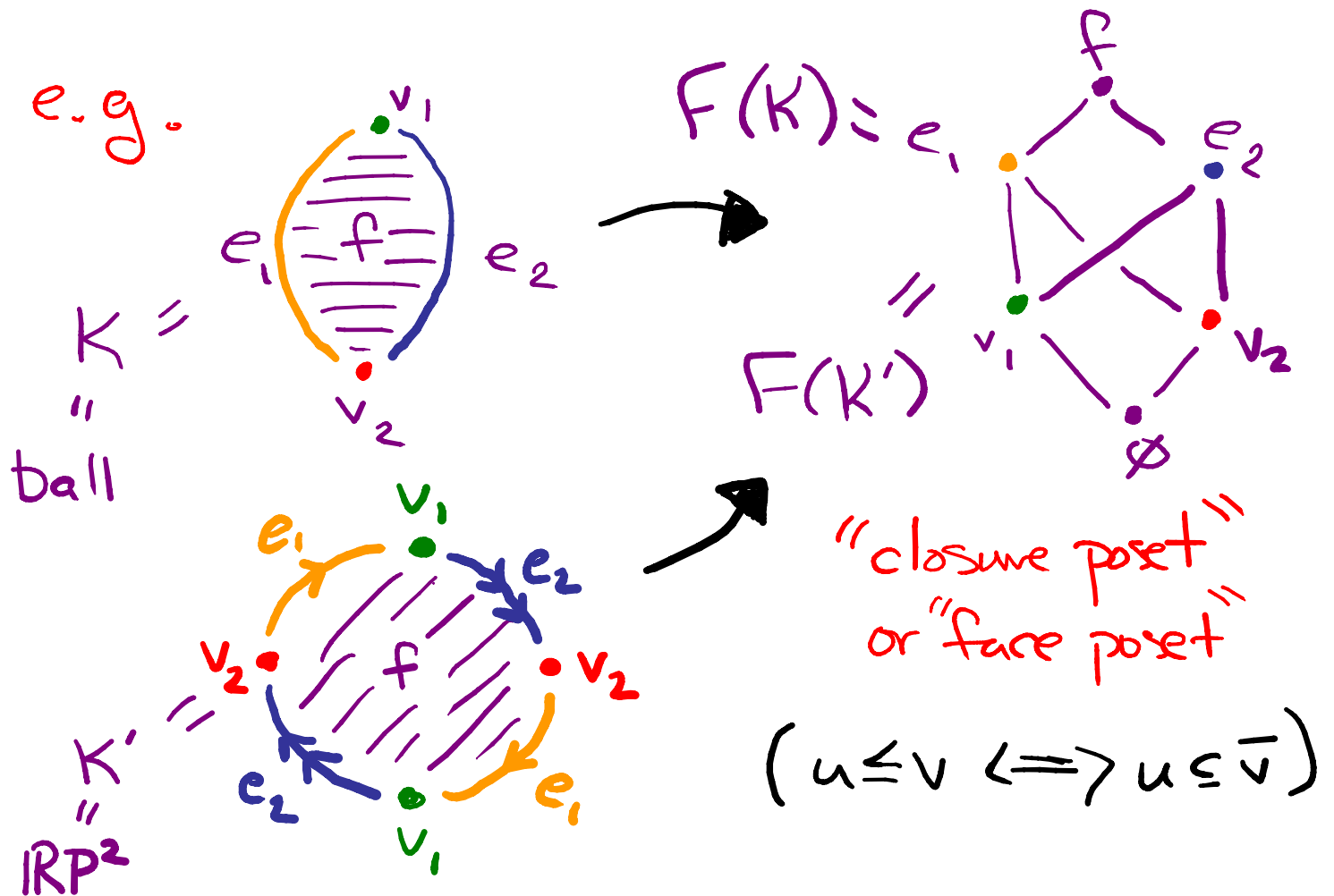
Shelling Face Posets of Stratified Spaces of Electrical Networks $\approx$ of Phylogenetic Trees

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-joint work with Rick Kenyon

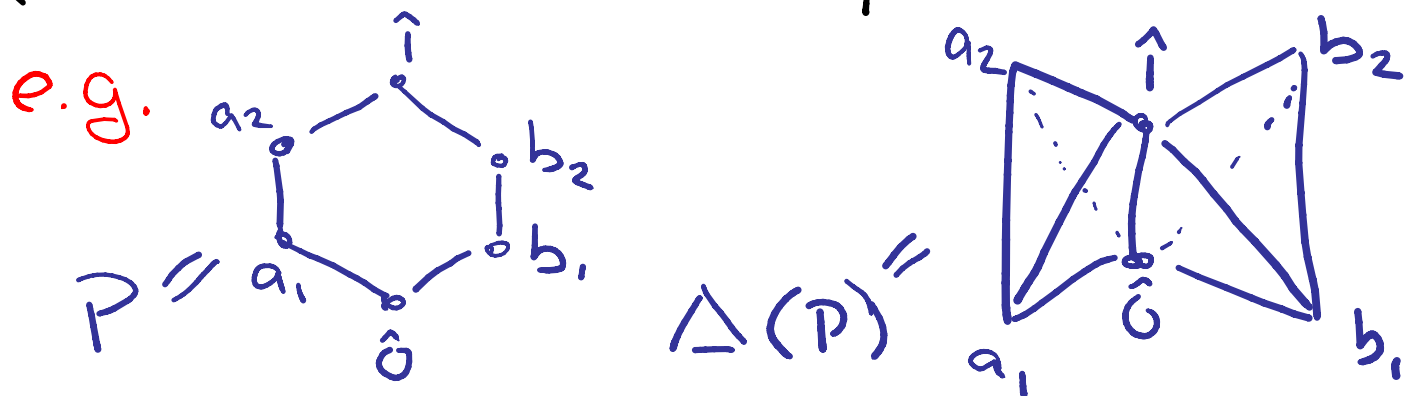
(slides at <https://pages.uoregon.edu/plherish/talks.html>)

CW Complexes $\neq$ their Face Posets



Recall: A **CW complex** is comprised of **open cells** each homeomorphic to an open ball. A **regular CW complex** further has cell closures homeomorphic to closed balls. e.g. **simplicial complexes**

Def'n: The **order complex** (or **nerve**) of a poset P is the abstract simplicial complex $\Delta(P)$ whose i -dimensional faces are the $(i+1)$ -"chains" $v_0 < \dots < v_i$ in P



Key Property (Hall; popularized by Rota):

$$M_P(x, y) = \tilde{\chi}(\Delta_P(x, y)) = \begin{matrix} -1 + \# \text{vertices} \\ - \# \text{edges} \\ + \# 2\text{-faces} \dots \end{matrix}$$

$$= -1 + \beta_0 - \beta_1 + \beta_2 - \dots$$

- $(u, v) = \{z \in P \mid u < z < v\}$
- $u < \cdot v$ means $u < v$ & $\nexists z$ s.t. $u < z < v$
- saturated chains u to $v := u < \dots < v$

Background on Face Posets & CW Posets

- A graded poset with $\hat{0} \neq \hat{1}$ is **Eulerian** if $\mu(u, v) = (-1)^{\text{rk}(v) - \text{rk}(u)}$ for all $u \leq v$.
- A graded poset P is a **CW poset** if
 - (1) $\hat{0} \in P$
 - (2) P has at least one other element
 - (3) $\Delta(\hat{0}, u) \cong S^{\text{rk}(u) - 2}$ for $u \neq \hat{0}$

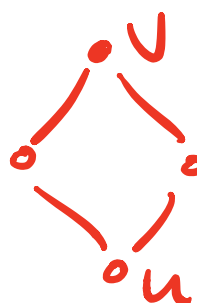
Thm (Björner): P is CW poset $\Leftrightarrow$

there exists regular CW complex
with P as poset of closure relns

Cor: CW Poset $\Rightarrow$ Eulerian

Some CW Posets

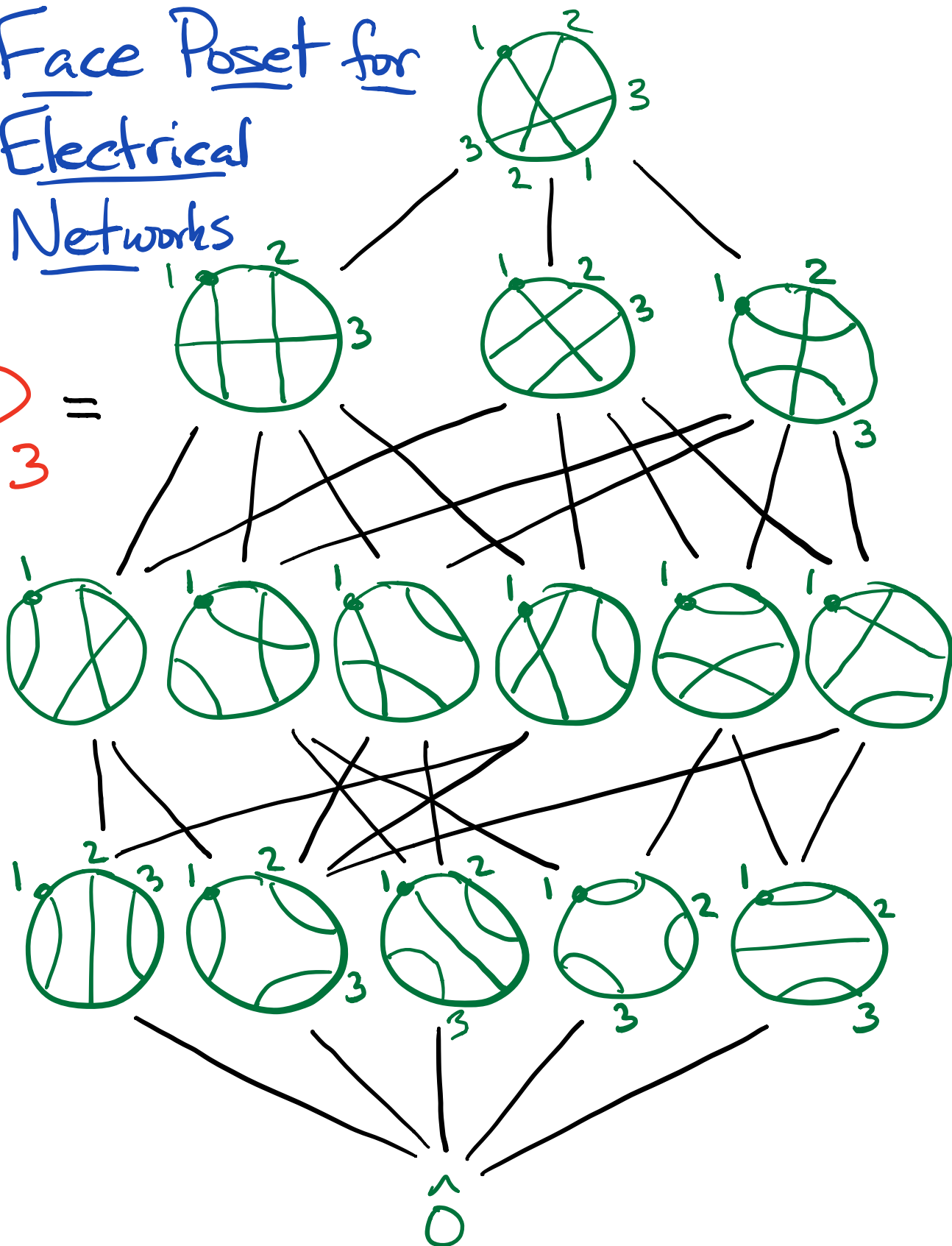
- all graded, thin, shellable posets (Danaraj-Klee)

"thin" $\Leftrightarrow$  $rk(v) - rk(u) = 2$
then $|[u, v]| = 4$

- Bruhat order (Björner-Wachs; Dyer)
- Face posets of stratified spaces of electrical networks
 - conjectured by Thomas Lam
 - proved by H.-Kenyon (a main topic for today's talk)

Face Poset for Electrical Networks

$P_3 =$



The Uncrossing Poset (Face Poset for Electrical Networks)

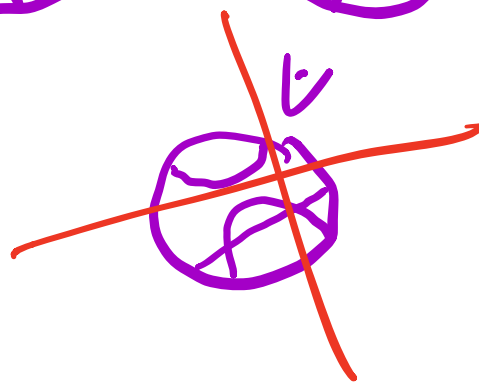
- $\hat{1} :=$ wire diagram w/ all $\binom{n}{2}$ crossings of n wires



- $u < v$ if u obtained from v uncrossing pair of wires without introducing double crossing



- $\hat{0}$ adjoined below Catalan many atoms



- $\mathcal{P}_n :=$ uncrossing poset with n wires

A Conjecture of Thomas Lam

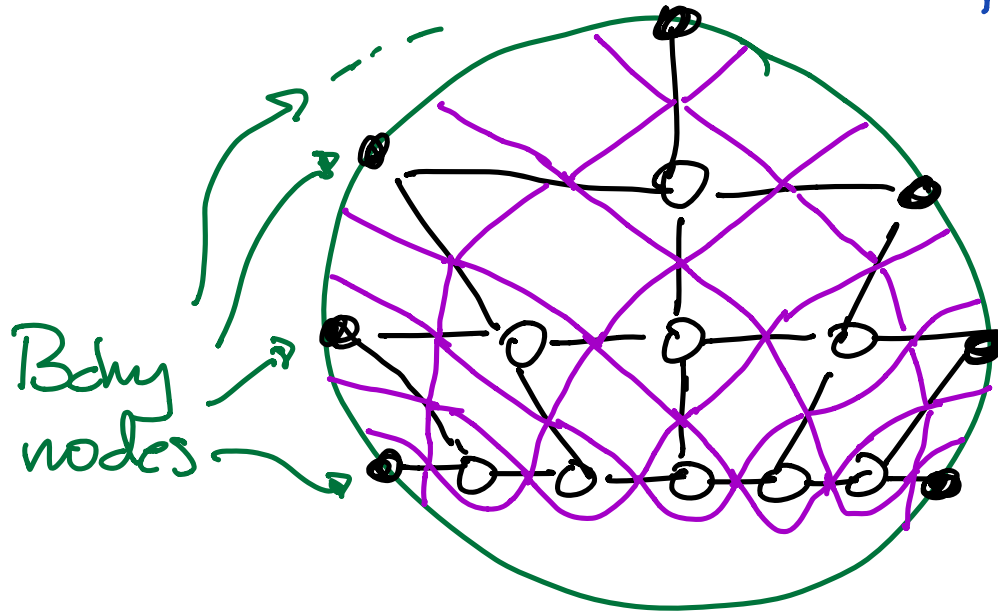
Thm (Lam): The uncrossing poset is Eulerian.

Conjecture (Lam): The uncrossing poset is lexicographically shellable.

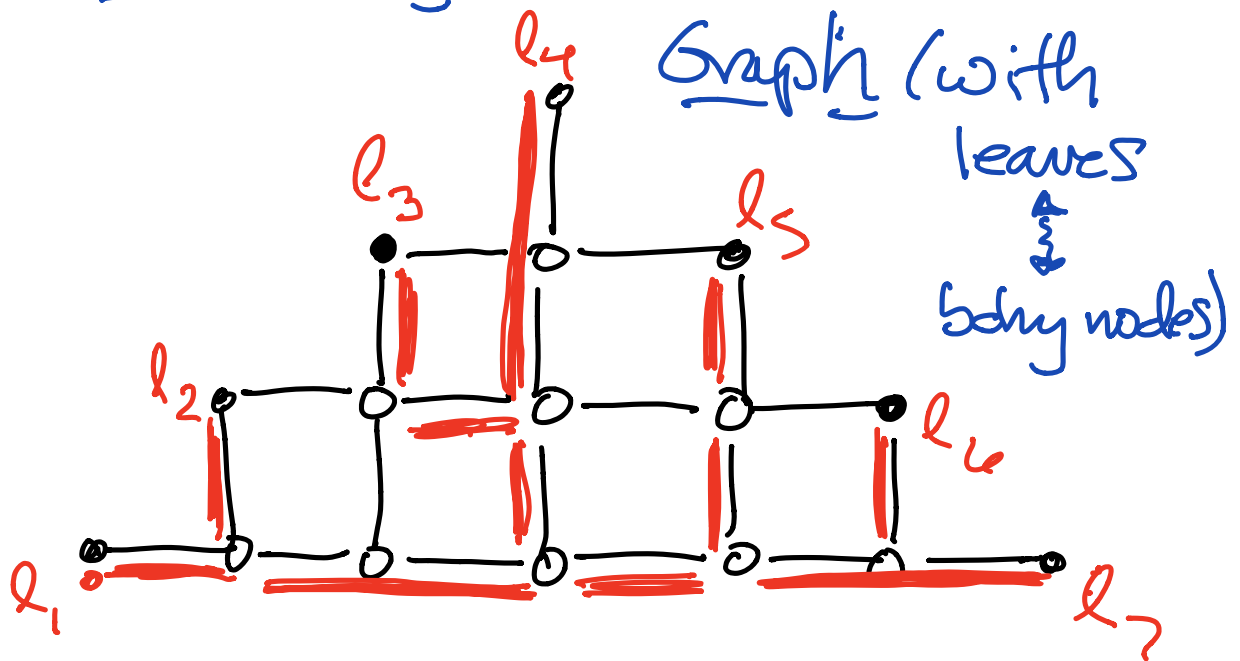
Thm (H.-Kenyon): Uncrossing posets are dual EC-shellable.

Cor: They are CW posets.

Dual Graph for $\hat{1} \in P_7$



Embedding Tree in Such Dual Graph (with



Edge Product Space, E_n , of Phylogenetic Trees with n leaves

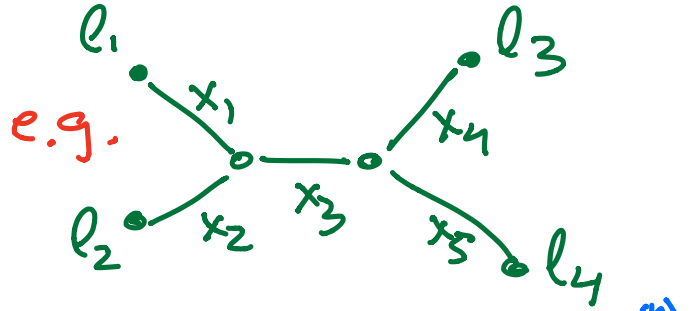
Step 1: Given tree T with (independent)

edge probabilities

$x_1, x_2, \dots, x_m$

leaves $l_1, l_2, \dots, l_n$,

may calculate vector $m_T(x_1, \dots, x_m) \in [0, 1]^{\binom{n}{2}}$
of probabilities for each pair of leaves that
path between them is fully present



e.g. $(\underbrace{x_1 x_2}_{P_{12}}, \underbrace{x_1 x_3 x_4}_{P_{13}}, \underbrace{x_1 x_3 x_5}_{P_{14}}, \underbrace{x_2 x_3 x_4}_{P_{23}}, \underbrace{x_2 x_3 x_5}_{P_{24}}, \underbrace{x_4 x_5}_{P_{34}})$

Step 2: Define edge-product space E_n
for trees with n leaves as

$$\bigcup_{T \text{ trees}} \text{im}(m_T) \subseteq [0, 1]^{\binom{n}{2}}$$

w/ leaves $\{l_1, \dots, l_n\}$ & nonleaves all $\deg \geq 3$

Theorem (Gill-Limussan-Moulton-Steel):

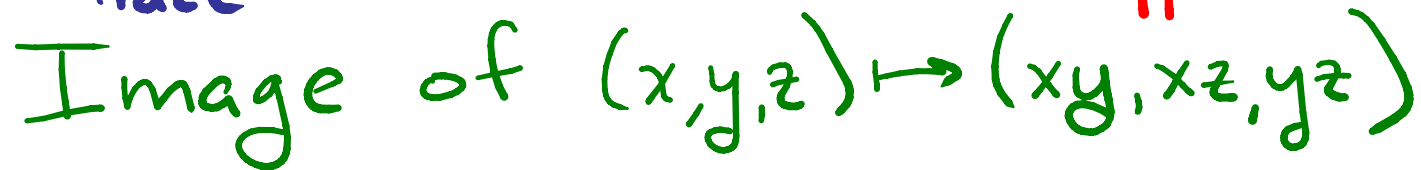
There exists a shelling for each interval in the face poset for E_n , namely the Tuffley poset.

Corollary (Gill-Limussan-Moulton-Steel):

The edge product space of phylogenetic trees is a regular CW complex,

Theorem (Stadnyk): The Tuffley poset in its entirety is not shellable.

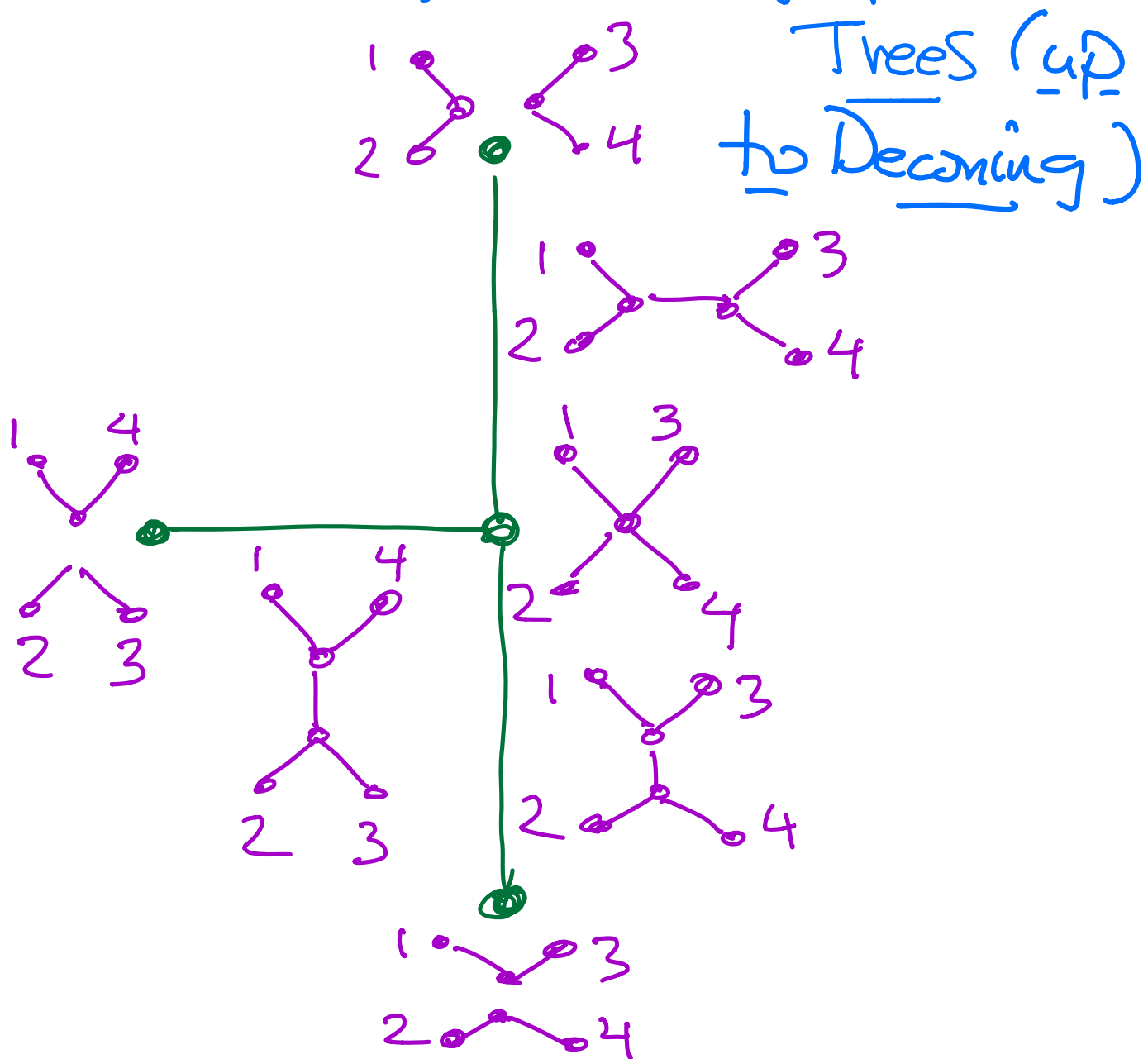
Theorem (H.-Kenyon): Each interval in the Tuffley poset is isomorphic to an interval in the uncrossing poset, hence has an explicitly constructed shelling by virtue of the uncrossing poset shelling.



(bounding surfaces $m_1, m_2 \leq m_3$
 $m_1, m_3 \leq m_2$ in $[0, 1]^3$
 $m_2, m_3 \leq m_1$)

Small Example of Edge

Product Space of Phylog.

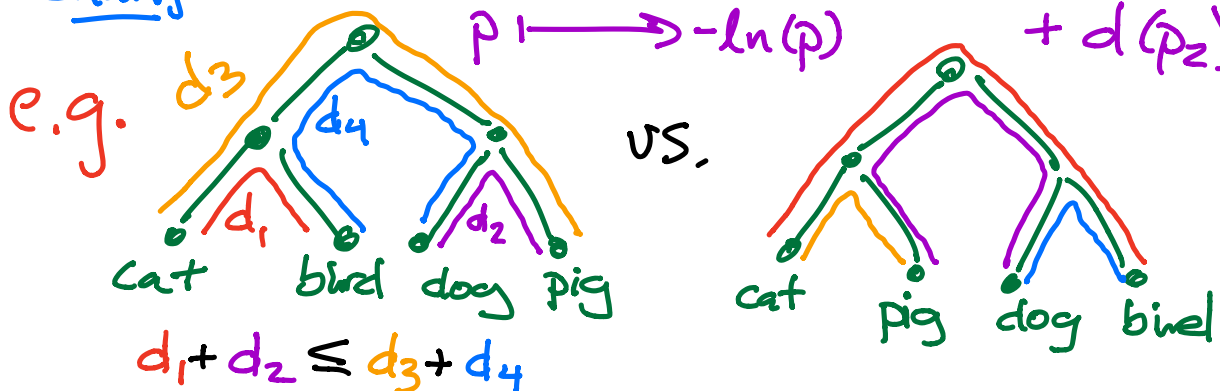


Motivations for Edge Product Space of Phylogenetic Trees

(1) Determining edge lengths
from distances between leaves
(data measurable from DNA)

(2) Determining tree type (i.e.
which cell tree lives in)

Convention: $(0, 1] \xrightarrow[\text{homeom}]{d} \mathbb{R}_{\geq 0}$ $P_1, P_2 \mapsto d(p_1)$
Change $p_1 \mapsto -\ln(p)$ $+ d(p_2)$



(1) + (2) $\Rightarrow$ likely phylogenetic
history from "distances"

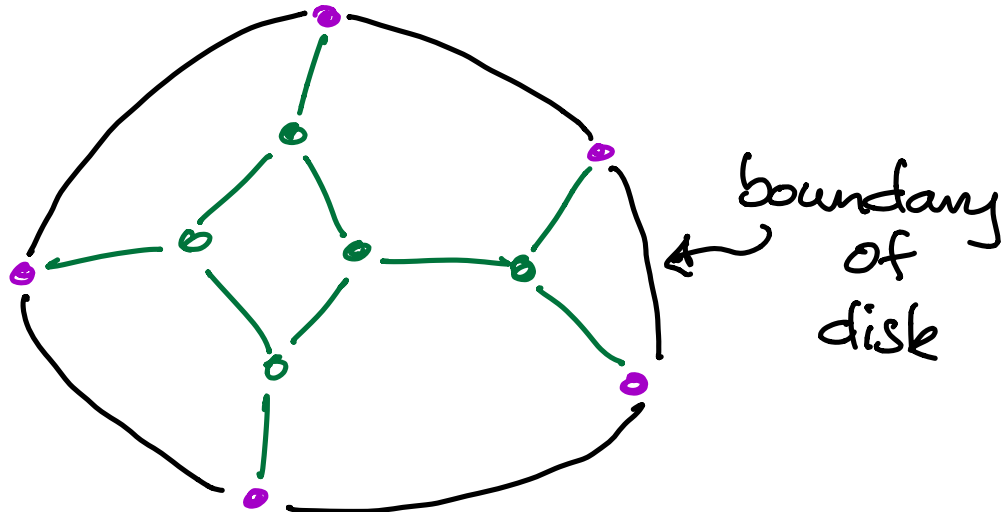
Some Earlier Work on E_n

- Billera-Holmes-Vogtmann: uniqueness of geodesics
- Owen-Provan: efficient algorithm to calculate geodesics
- Moulton-Steel: CW complex
- Ardila-Klivans: homeomorphism (via subdivision) to order complex of partition lattice
- Basu-Gabrielov-Vorobiy: "toric cubes" are regular CW complexes

Note: E_n is compactification of "tree space" in BHV, OP, AK

Space of "Response Matrices" of a Planar Electrical Network

e.g.



I = interior
nodes

N = boundary
nodes

$$\underbrace{\begin{pmatrix} A & B \\ B^T & C \end{pmatrix}}_{\substack{\text{vector} \\ \text{of} \\ \text{voltages}}} \underbrace{\begin{pmatrix} V_N \\ V_I \end{pmatrix}}_{\substack{\text{vector} \\ \text{of} \\ \text{currents}}} = \underbrace{\begin{pmatrix} C_N \\ 0 \end{pmatrix}}_{\substack{\text{vector} \\ \text{of} \\ \text{currents}}}$$

$$\underbrace{(A - BC^{-1}B^T)}_{\text{response matrix}} V_N = C_N$$

"response matrix" of the network

A Goal: Given a graph G , study the space of response matrices

as image of

$$f: \left\{ \begin{array}{c} \text{conductance} \\ \text{vectors} \end{array} \right\} \rightarrow \left\{ \begin{array}{c} \text{response} \\ \text{matrices} \end{array} \right\}$$

$$(\mathbb{R}_{\geq 0} \cup \{\infty\})^{|E|}$$

Note:

contracting an edge $\longleftrightarrow$ sending conductance to ∞ (i.e. resistance to 0)

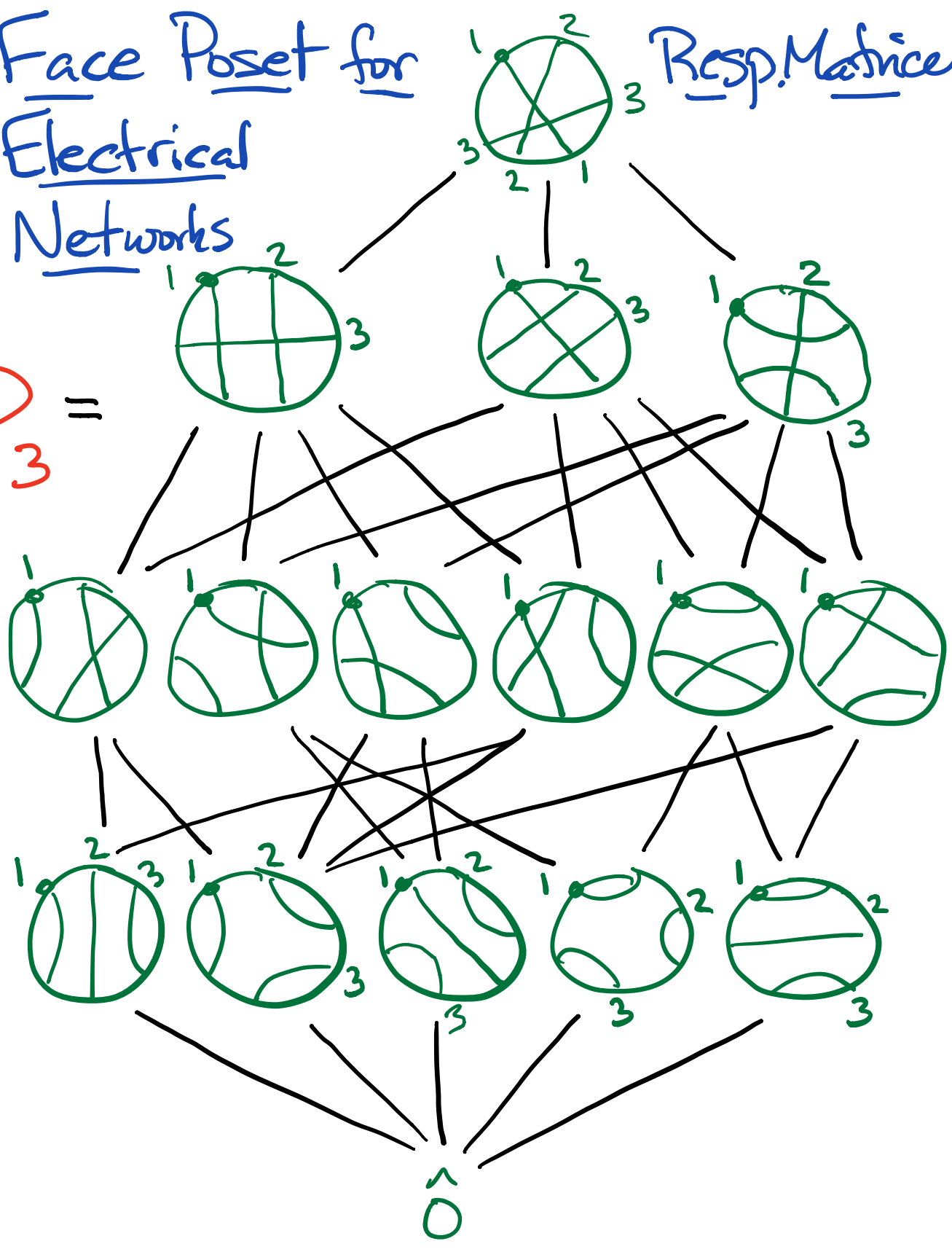
deleting an edge $\longleftrightarrow$ sending conductance to 0 (resistance to ∞)

Secondary Goal: Study fibers of f

Face Poset for Electrical Networks

Resp. Matrices for

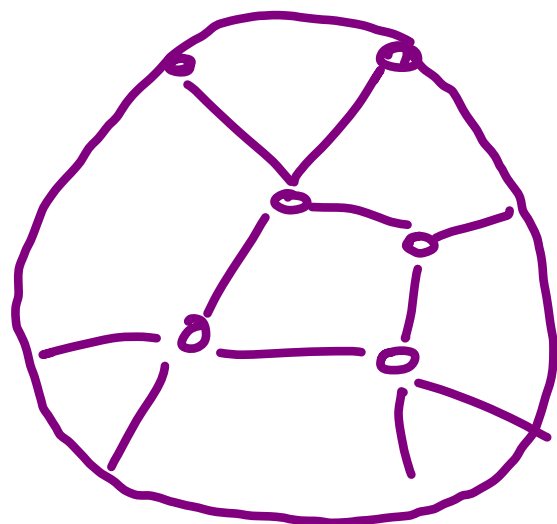
$P_3 =$



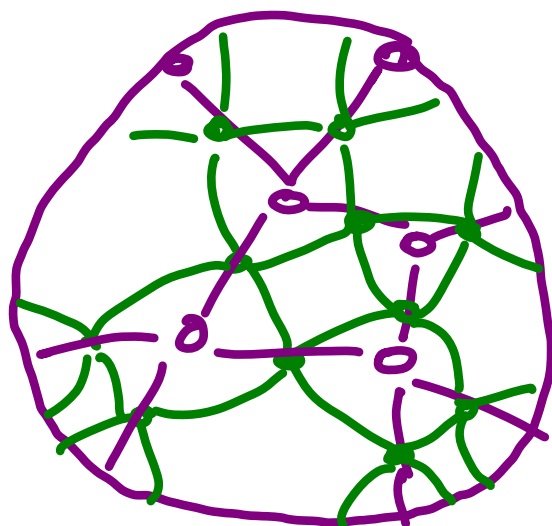
Correspondence: From Graphs to

Uncrossing
Wire Diagrams
"Medial Graphs"

$\Gamma =$

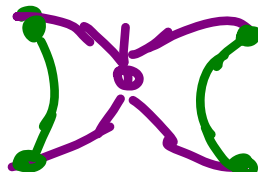
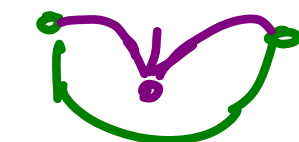
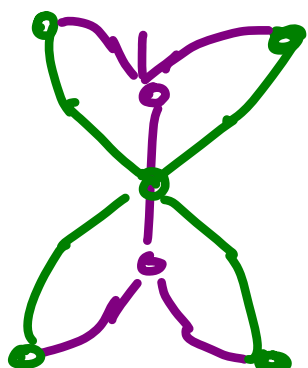


$G(\Gamma) =$



= wire
diagram

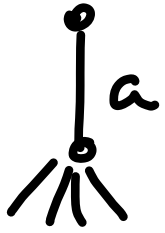
Uncrossing:

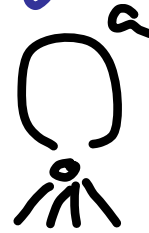


deletion

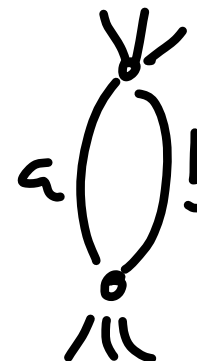
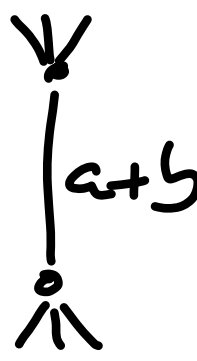
contraction

Turning to Fibers via "Electrical Equivalence"

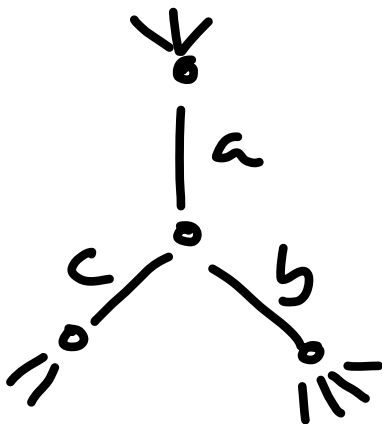
(1)  $\sim$ 

(2)  $\sim$ 

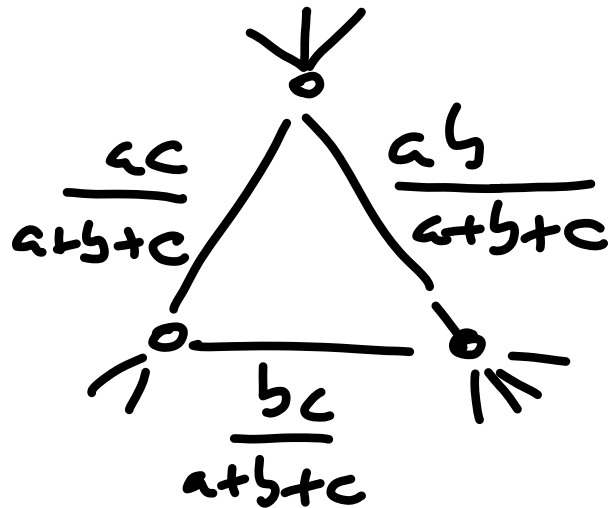
(3)  $\sim$ 

(4)  $\sim$ 

(5) "Y- Δ moves"

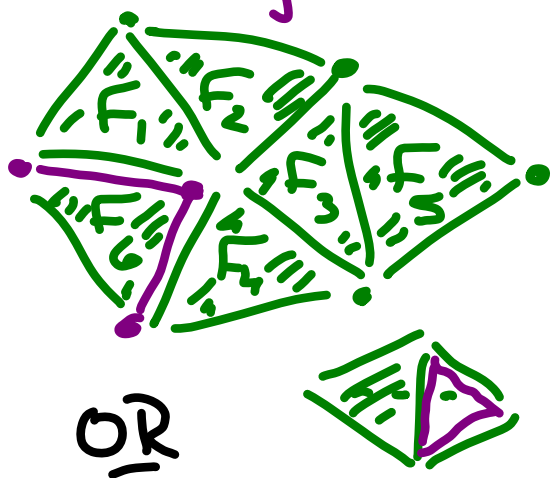


$\sim$



Shellability

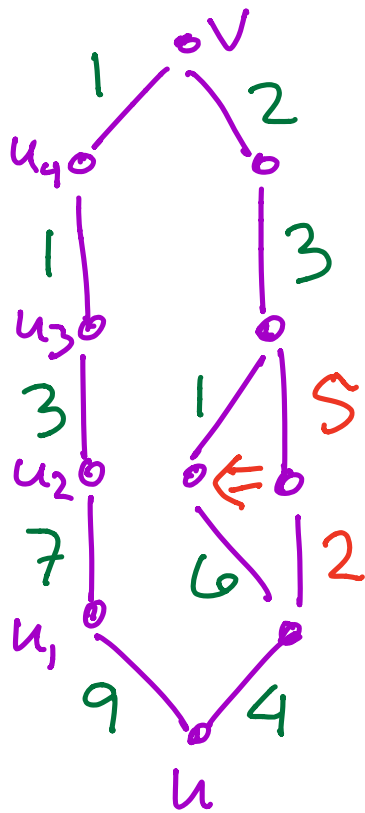
- Simplicial complex is **pure** of dim. d if all maximal faces ("facets") are d -dimensional
- simplicial complex is **shellable** if there is total order $F_1, F_2, \dots, F_k$, a **shelling**, on facets s.t. $\bar{F}_j \cap (\bigcup_{i < j} \bar{F}_i)$ is pure, codimension one subcomplex of $\bar{F}_j$ for each $j > 1$ (hence is $\partial \bar{F}_j$ or has a cone point).



- Each facet attachment preserves homotopy type or closes off a new sphere

Lexicographic Shellability for (Björner, Björner-Wachs) $\Delta(P)$

A poset is **dual EL-shellable** if it admits edge labeling λ (called **dual EL-labeling**) of cover relations $x \lessdot y$ s.t. for each $u < v$ in P :



(1) $\exists!$ Saturated chain

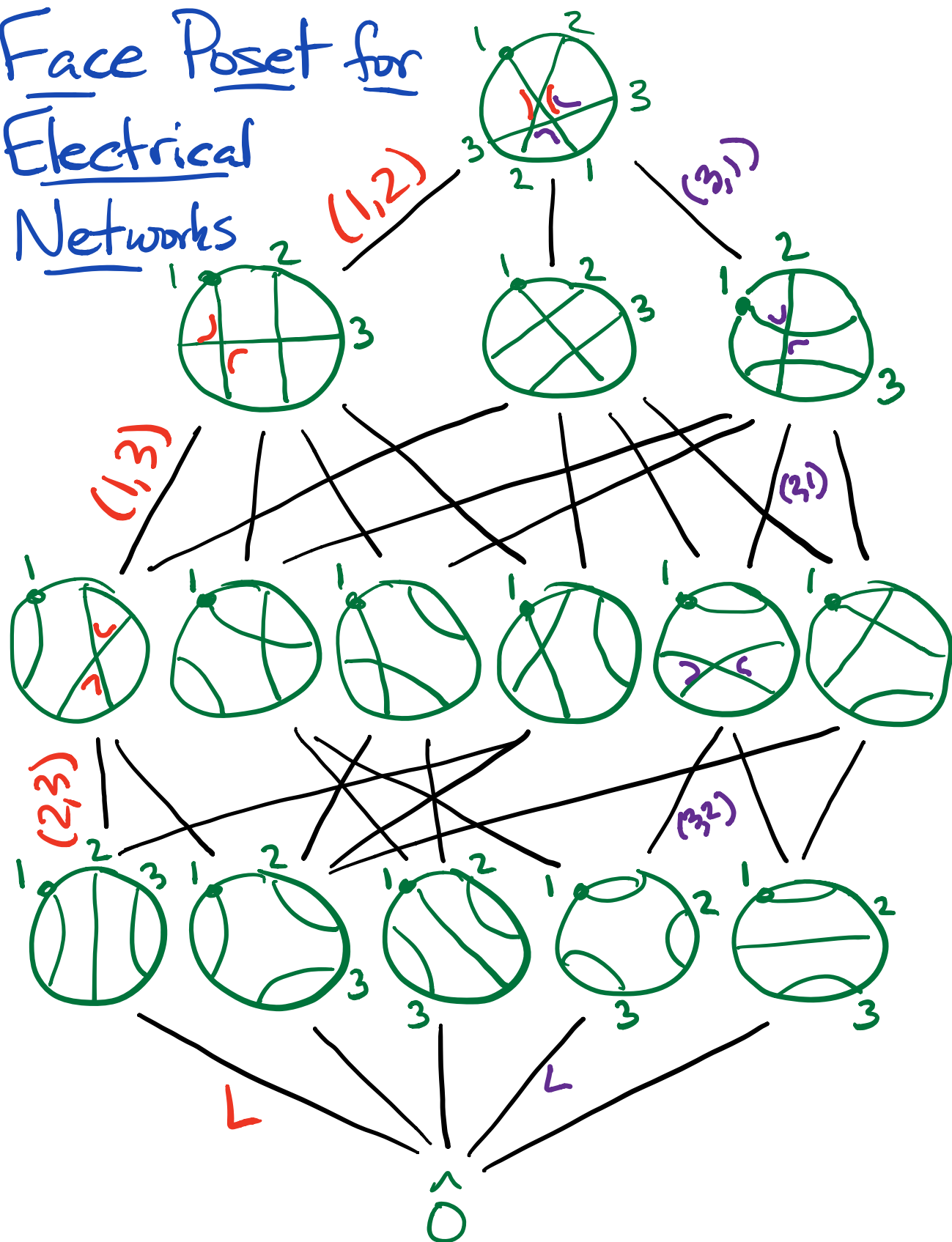
$$u < u_1 < u_2 < \dots < u_k < v \text{ s.t.}$$

$$\lambda(u_k, v) \leq \lambda(u_{k-1}, u_k) \leq \dots \leq \lambda(u_1, u_1)$$

and

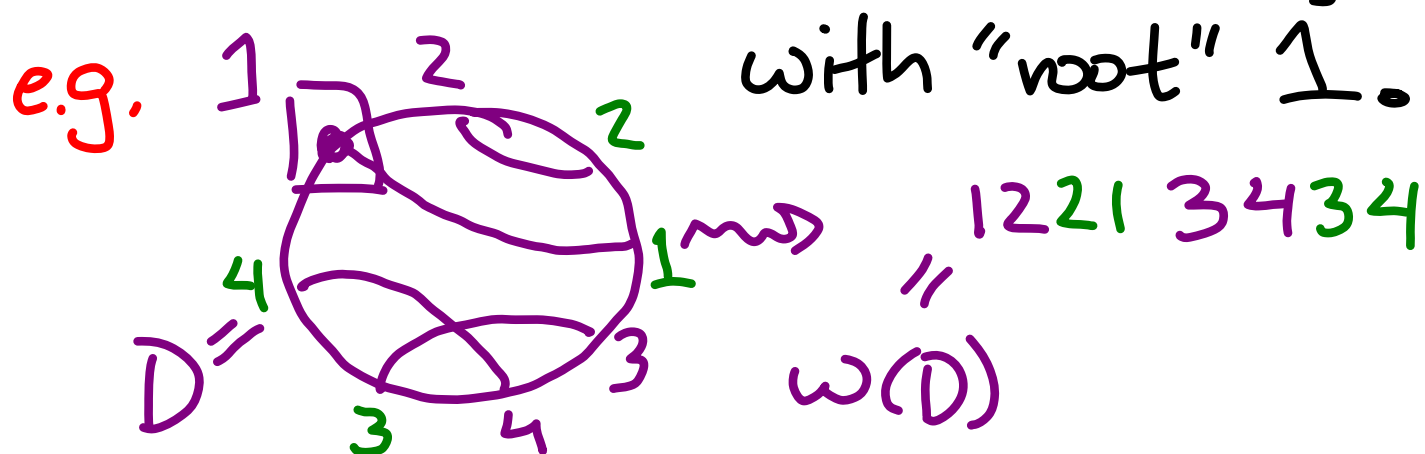
(2) $(\lambda(u_k, v), \dots, \lambda(u_1, u_1))$ is smaller in dictionary order than all other label sequences from v to u

Face Poset for Electrical Networks



Edge Labeling

Step 1: Define **word** of wire diagram D , denoted $w(D)$, as sequence of $2n$ wire endpoints encountered clockwise starting



Step 2: Label $D \leftarrow D'$ for $i < j$ as

- (i, j) if $ijij$ in $w(D')$ becomes $ijji$ in $w(D)$
- (ji) if $ijij$ becomes $ijij$

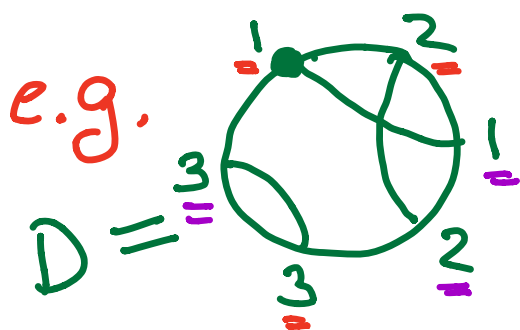
Label Ordering

- labels $\{(i,j) \mid 1 \leq i < j \leq n\} \cup \{L\}$
 $\cup \{(j,i) \mid 1 \leq i < j \leq n\}$
- $(i,j) < L < (r,s)$ for all $i < j$
and all $r > s$
- $(1,2) < (1,3) < (1,4) < \dots < (1,n) < (2,3)$
 $< (2,4) < \dots < (2,n) < (3,4) < \dots < (3,n)$
 $< \dots < (n-1,n)$
- $(n,n-1) < (n,n-2) < (n-1,n-2) < (n,n-3)$
 $< (n-1,n-3) < (n-2,n-3) < \dots$
 $< (2,1)$

Rk: finite type A reflection order $\{(i,j)\}$
then L , then reversal for $\{(j,i)\}$

"Start Sets" and Connection to Type A Bruhat Order

The start set of D , denoted $S(D)$, is the subset of $\{1, 2, \dots, 2n\}$ of positions in $w(D)$ where 1st copies of letters occur.



$$w(D) = \underline{1} \underline{2} 1 2 \underline{3} 3$$

$$S(D) = \{1, 2, 5\}$$

$$\pi(D) = 123 \in S_3$$

Propn: If $D' < D$ and $S(D') = S(D)$,

then $[D', D] \cong [\underbrace{\pi(D'), \pi(D)}_{\text{subwords of } w(D) \neq w(D)}]_{\text{Bruhat}}$

- uncrossing order labeling coincides here w/ Dyer's Bruhat order labeling

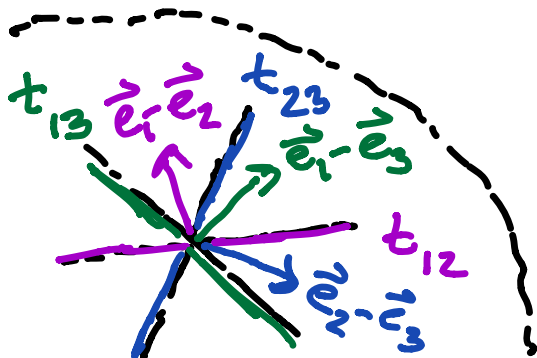
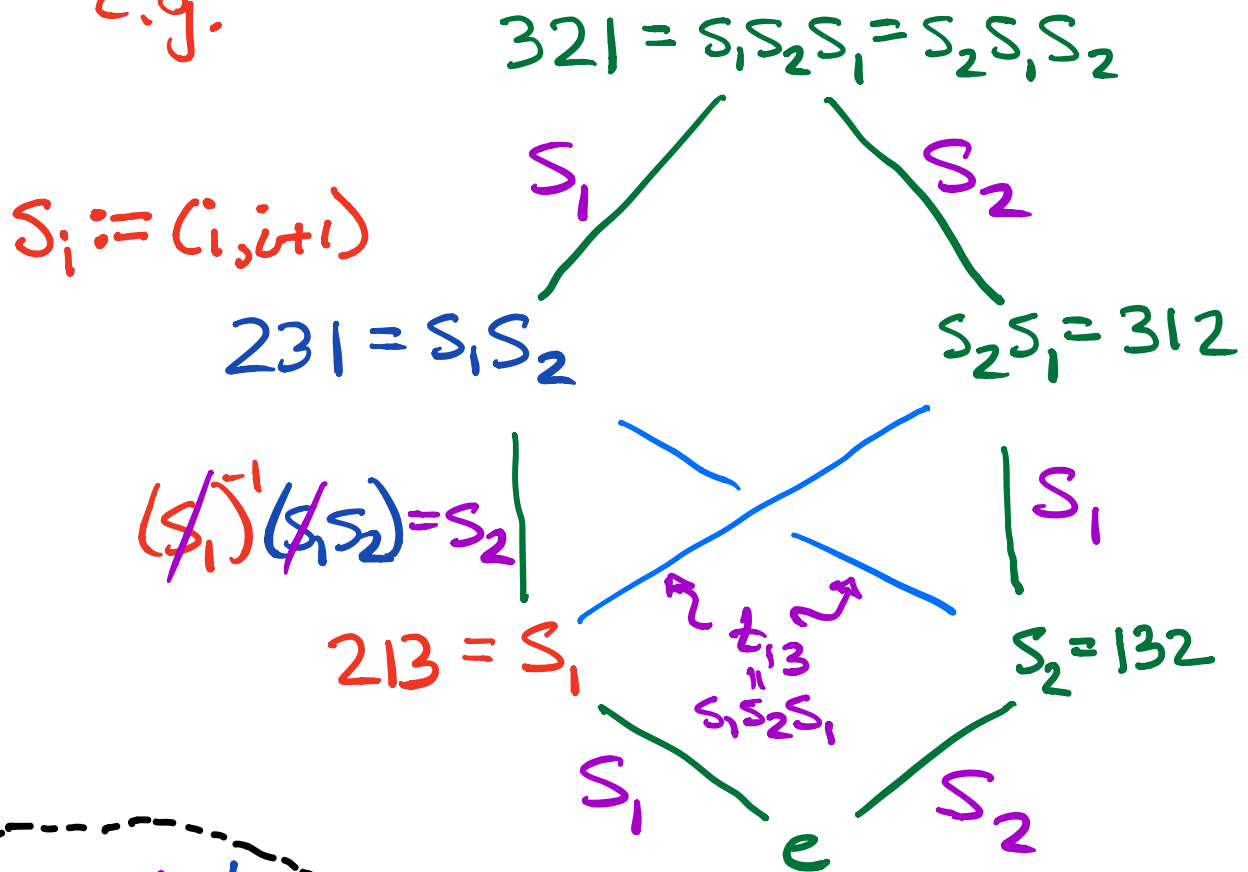
Concluding Remarks

1. Galashin-Karp-Lam recently announced proof that stratified space of response matrices is regular CW complex.
2. Analogy w/ my past work tot. nonneg. part unip. radical of Borel?
3. Work in progress: graphical analogue of subword complexes to describe fibers.

Thanks for your attention!

Bruhat Order of a (Finite) Reflection Group / Coxeter Group $\hat{=}$ Dyer's EL-Labeling

e.g.



EL-labeling: $u \prec v = ut$

$$\lambda(u, v) := u^{-1}v = t$$

Dyer's EL-labeling (cont)

- Use any "reflection order" to totally order edge labels
- Dyer proved these exist & induce EL-labelings

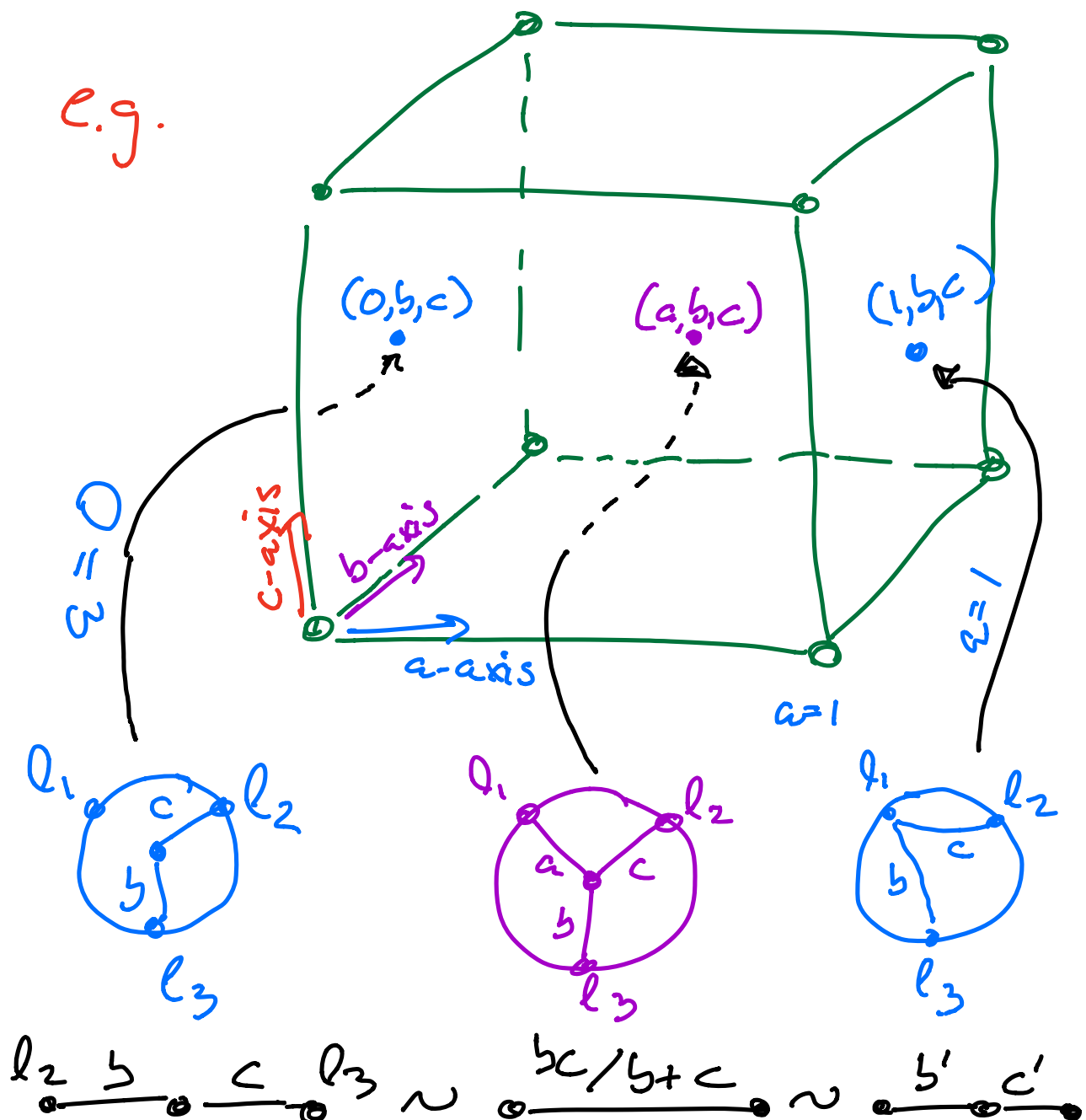
Defn: A total order on positive roots ($\neq$ assoc'd reflections) is **reflection order** if $\alpha < c_1\alpha + c_2\beta < \beta$ or $\beta < c_1\alpha + c_2\beta < \alpha$ for each such triple of positive roots w/ $c_1, c_2 > 0$

e.g $(1,2) < (1,3) < (2,3)$ or

$(2,3) < (1,3) < (1,2)$ in type A

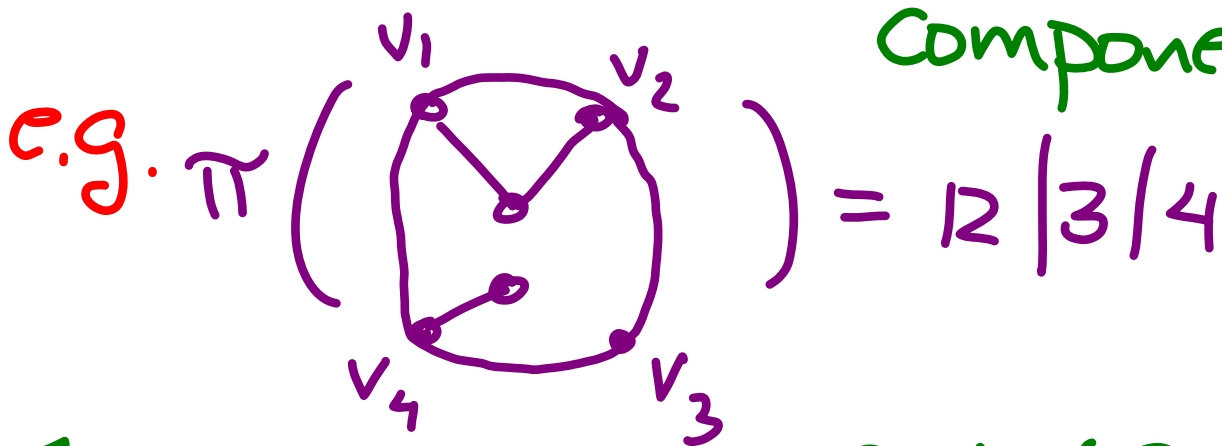
Stratification into "Cells" for Space of Response Matrices

e.g.



Minors of Response Matrix via "Gours"

$\pi(\tilde{G}) :=$ set partition of bdy
graph nodes into connected
components



Thm (Special case of Next Result):

$$L_{ij}(G) = \frac{\sum_{\substack{G' \leq G \\ \pi(G') = ij | \text{singletons}}} \text{wt}(G')}{\sum_{\substack{G' \leq G \\ \pi(G') = \text{all singletons}}} \text{wt}(G')}$$

$\text{wt}(G') =$ product
of
edge
weights
(i.e. conductances)

$\sum_{\substack{G' \leq G \\ \pi(G') = \text{all singletons}}} \text{wt}(G')$

Lam's Dual Embedding of Uncrossing Order into $\tilde{S}_{2n}$ Braid Order

• $D_e = \text{fully crossed diagram}$ $\rightsquigarrow \exists! D \in \tilde{S}_{2n}$
 $\exists \gamma \text{ s.t.}$
 $i \rightarrow i+n$

• $D < D'$ where D obtained from D' by i, j wire endpoint swap $\rightsquigarrow$
 $\exists D$
 $\parallel$
 $(i, j) g_{D'} (i, j)$



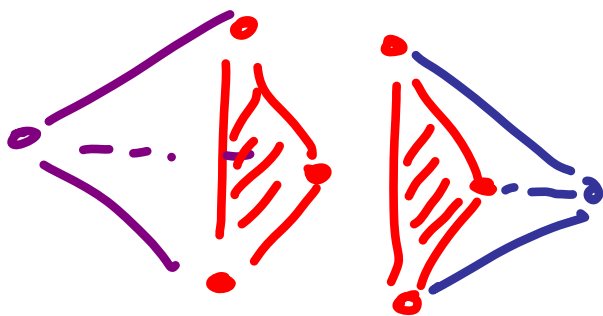
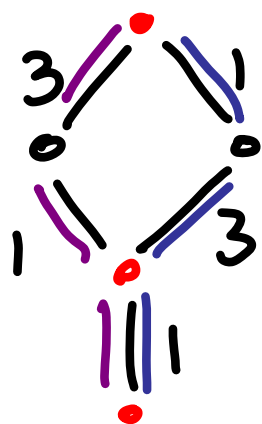
$$g_D = (1, 2) g_{D'} (1, 2)$$

$$= (4, 5) g_{D'} (4, 5)$$

Thm (Björner): EL-labeling $\Rightarrow$ Shelling

Idea: Lexicographic order on maximal chains (breaking ties arbitrarily) induces shelling order on corresponding facets of $\Delta(P)$.

- "descents in labeling" $\iff$ codim. one overlap of facets



- "descending" chains $\iff$ facets attaching along entire boundary $\iff$ spheres

- $M_P(u, v) = \pm \# \text{descending chains } u \text{ to } v$
(for P graded)